

Mathematical Model of Taxi Routing

Goal of taxi routing:

Determine the minimum number of taxis needed to serve all the trips without causing delays (with 20/20 hindsight)

Input to the model:

- A collection of *n* taxi trips
- $N = \{(p_1, T_1^p, d_1, T_1^d), \dots, (p_n, T_n^p, d_n, T_n^d)\}, \text{ where the } i^{th} \text{ trip}$ picks up a passenger at time T_i^p at point p_i and drops off the fare at time T_i^d at point d_i , $i \in \{1, \dots, n\}$
- Maximum allowable taxi waiting time δ (in minutes)

Construct a bipartite graph G = (D, P, E):

- $D = \{(d_1, T_1^d), \dots, (d_n, T_n^d)\}$ denotes the set of drop-off nodes
- $P = \{(p_1, T_1^p), \dots, (p_n, T_n^p)\}$ denotes the set of pickup nodes
- $E = \left\{ \left\{ (d_i, T_i^d), (p_j, T_j^p) \right\} : T_j^p T_i^d \delta \le time(d_i, p_j) \le T_j^p T_i^d \right\}$
- Elapsed time between the drop-off time for d_i and the pickup time for p_j , that is, $T_i^p - T_i^d$, is at least the time needed to travel between these two points $time(d_i, p_i)$ and at most $time(d_i, p_i) + \delta$; hence, the taxi can reach the new pickup in time, and does not need to wait more than δ minutes for the pickup to be ready

Lemma Given any matching of size m, there exists a covering of all *n* trips with n - m taxis

• Application: Maximum cardinality matching gives rise to the minimum number of taxis that cover all the trips [1][2]



- An example of 10 trips that can be covered by 4 taxis
- \Box Maximum cardinality matching (m = 6) is colored in red
- Nodes are colored with respect to their corresponding taxi routes

Taxi Routing Optimization

Tinghan (Joe) Ye Advisor: David Shmoys School of Operations Research and Information Engineering, Cornell University

Computational Results

The optimal taxi routing (with $\delta = 10$) is compared with the actual one

- Total number of taxis that covers all trips reduces from 11568 to 7475, a 35% reduction
- Total trip time and empty trip time are greatly reduced, from 16.1 to 4.2 hours and from 14.3 hours to 1.4 hours, respectively, on average
- To offset the reductions in trip time, the number of trips and on-trip percentage are greatly increased, from 11.3 to 17.5 and from 17.1% to 72.6% respectively, on average



A comparison of the time distribution of the optimal taxi fleets (with breakdown) and the actual taxi fleets • On average, number of circulating taxis reduces from 7748 to 1300, an 83% reduction



References

[1] M. M. Vazifeh, P. Santi, G. Resta, S. H. Strogatz & C. Ratti, "Addressing the minimum fleet problem in on-demand urban mobility." Nature 557.7706 (2018): 534-538.

[2] X. Zhan, X. Qian & S. Ukkusuri, "A graph-based approach to measuring the efficiency of an urban taxi service system." IEEE Transactions on Intelligent *Transportation Systems* 17.9 (2016): 2479-2489.

Min-max theorem is an important concept in solvable discrete optimization problems; we prove a min-max theorem for the taxi routing problem

• Consider a simplified setup:

Definition Any two trips N_i and N_j are considered compatible with each other if a taxi can reach from d_i to p_j or from d_j to p_i in time

 $|K| = k \text{ and } S = \bigcup_{i \in K} \{p_i, d_i\} \subseteq I$

Lower bound An independent set of size n+k (nodes in shaded region) leads to k trips that are pairwise incompatible (orange box)

Theorem

The maximum size of a set of trips that are pairwise *incompatible* is equal to the minimum number of taxis needed to cover all trips

Proof.



waiting time for the new pickup

Theoretical Results

 $E = \{ \{ (d_i, T_i^d), (p_j, T_j^p) \} : time(d_i, p_j) \le T_j^p - T_i^d \};$ i.e., no upper bound on taxi waiting time ($\delta = \infty$)

Proposition Consider two disjoint sets $P = \{p_1, \dots, p_n\}$ and $D = \{p_1, \dots, p_n\}$ $\{d_1, \dots, d_n\}$ and a subset $I \subseteq P \cup D$, where |I| = n + k, for some $k \ge 0$; then there exists a subset $K \subseteq \{1, \dots, n\}$ where

certificate:



Future Work

• The dual object when there's an upper bound (δ) to the

• An optimal taxi routing that accounts for secondary objectives of waiting time (e.g., minimize maximum waiting time)